# APPLICATION OF GRAPH THEORY IN NETWORK ANALYSIS 

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#### Abstract

Graph theory provides the basis for many network analysis techniques. Basic graph theory concepts are very general and can be applied to a wide variety of network problems such as topological design, routing reliability analysis, and network capacity. Specific examples demonstrate that graph theory is a practical tool for solving network and distributed system problems.


Keywords: Adjacency matrix, Incidence matrix, reliability.

## INTRODUCTION

Any network problem can be represented by graph and graph can be represented by two different ways inside a computer, namely by using the adjacency matrix or the incidence matrix of a graph.

1 Adjacency matrix: If a graph $G$ has $n$ vertices,listed as $\mathrm{v}_{1}, \mathrm{~V}_{2}, \ldots, \mathrm{v}_{\mathrm{n}}$. The adjacency matrix of G , with respect to this particular listing of the $n$ vertices of $G$,is the $n \times n$ matrix $A(G)=\left(A_{i j}\right)$ where the $(1, j)$ th entry $v_{i}$ to the vertex $v_{j}$. The following figure 1 shows a graph $G$ with vertices listed as $\mathrm{v}_{1}, \ldots, \mathrm{~V}_{4}$ and its adjacency matrix $\mathrm{A}(\mathrm{G})$ with respect to this listing. Figure 1:A graph and its adjacency matrix


A(G):4×4 matrix

|  | v1 | v2 | v3 | v4 |
| :---: | :---: | :---: | :---: | :---: |
| V1 | 1 | 1 | 0 | 1 |
| V2 | 1 | 0 | 1 | 0 |
| V3 | 0 | 1 | 0 | 2 |
| V4 | 1 | 0 | 2 | 0 |

Here in $A(G)$ we have $\mathrm{a}_{\mathrm{ij}} \mathrm{a}_{\mathrm{ji}}$ for each I and j .Amatrix with this property is called symmetric.Also note that if $G$ has no loops then all the entries of the main diagonal of $A(G)$ are 0 , while if $G$ has no parallel edges then the entries of $A(G)$ are either 0 or 1.

2 Incidence matrix: Suppose that G has n vertices,listed as $\mathrm{v}_{1}, \mathrm{v}_{2}, \ldots, \mathrm{v}_{\mathrm{n}}$ and t edges,listed as $\mathrm{e}_{1}, \mathrm{e}_{2}, \ldots, \mathrm{e}_{\mathrm{t}}$. The incidence matrix of G , with respect to these particular listings of the vertices and edges of $G$,is the $n \times t$ matrix $M(G)=\left(m_{i j}\right)$ where the vertex $v_{i}$ is incidence with the edge $e_{j}$ i.e,
$M_{i j}=\left\{\begin{array}{l}0 \text { if } v_{i} \text { is not end of } e_{j} \\ 1 \text { if } v_{i} \text { is an end of the non-loop } e_{j} \\ 2 \text { if } v_{i} \text { is an end of the loop } e_{j}\end{array}\right.$

The following figure 2 shows a graph $G$, with four vertices $v_{1}, \ldots, v_{4}$ and six edges $e_{1}, \ldots, e_{6}$ and its incidence matrix $M(G)$ with respect to these listings of the vertices and edges.

Figure 2: A graph and its incidence matrix

$\mathrm{M}(\mathrm{G})$ : A $4 \times 6$ matrix

|  | e1 | e2 | e3 | e4 | e5 | e6 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| v1 | 2 | 1 | 0 | 0 | 0 | 1 |
| v2 | 0 | 1 | 1 | 0 | 0 | 0 |
| v3 | 0 | 0 | 1 | 1 | 1 | 0 |
| v4 | 0 | 0 | 0 | 1 | 1 | 1 |

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