

Boolean Matrix Transpose Algorithm follows Parallel Computing Strategy

Sanil Shanker KP

Department of Computer Science, Farook College
Kozhikode, India

Abstract: In 2016, Sanil put forward an algorithm to transpose Zero- One matrix. This paper explains the algorithm's parallel nature with the support of an example.

Keywords: Boolean Matrix, Parallel computing.

I. INTRODUCTION

A matrix with entities that are either zero or one is called a Boolean matrix [1]. The transpose of a matrix A of order m x n is achieved by exchanging the rows and columns. In Parallel computing, there exists combination of operations to solve a single task [2]. This paper illustrates the parallel computing strategy of the proposed Boolean Matrix Transpose algorithm [3].

II. THE METHOD

Let **A** be a Boolean Matrix of order n x m. Then the transpose of **A** can be computed as:

1. Initialize the matrix A of order n x m .
2. Create the reference matrix .
3. Compute A^T by combining the characteristics of logical AND with logical OR operations in parallel way with the reference matrix.

III. PARALLEL STRATEGY

Consider the following example, Let **A** be a Boolean Matrix of order 4 x 5.

That is, **A** =>

0	1	1	0	1
0	0	0	0	0
1	0	0	1	0
0	0	0	0	0

A operates logical AND (\square) with reference matrix **D** as follows:

Reference Matrix **D**

\square

•	1	0	0	0
•	0	1	0	0
•	0	0	1	0
•	0	0	0	1

↓ ↓ ↓ ↓

z y x w

Here, the values filled in **z**, **y**, **x** and **w** are according to the formulation,

$$\sum_{i=1}^p W_{i,j}, \text{ where } j = 1, 2, \dots, q.$$

That is the first input binary vector operates logical AND with reference matrix **D** in parallel way as follows:

Reference Matrix D

1 (i=1)	•	1 (i=1,j=1)	0 (i=1,j=2)	0 (i=1,j=3)	0 (i=1,j=4)
0 (i=2)	•	0 (i=2,j=1)	1 (i=2,j=2)	0 (i=2,j=3)	0 (i=2,j=4)
0 (i=3)	•	0 (i=3,j=1)	0 (i=3,j=2)	1 (i=3,j=3)	0 (i=3,j=4)
0 (i=4)	•	0 (i=4,j=1)	0 (i=4,j=2)	0 (i=4,j=3)	1 (i=4,j=4)
		↓	↓	↓	↓
		z	y	x	w

(i) The value of **z**:

$$\sum_{i=1}^4 W_{i,j=1} = W_{1,1} + W_{2,1} + W_{3,1} + W_{4,1} = 1 + 0 + 0 + 0 = 1$$

Where,

$$W_{1,1} = I_{\text{Text}(i=1)} \cdot D_{i=1, j=1} = 1 \cdot 1 = 1$$

$$W_{2,1} = I_{\text{Text}(i=2)} \cdot D_{i=2, j=1} = 0 \cdot 0 = 0$$

$$W_{3,1} = I_{\text{Text}(i=3)} \cdot D_{i=3, j=1} = 0 \cdot 0 = 0$$

$$W_{4,1} = I_{\text{Text}(i=4)} \cdot D_{i=4, j=1} = 0 \cdot 0 = 0$$

(ii) The value of **y**:

$$\sum_{i=1}^4 W_{i,j=2} = W_{1,2} + W_{2,2} + W_{3,2} + W_{4,2} = 0 + 0 + 0 + 0 = 0$$

Where,

$$W_{1,2} = I_{\text{Text}(i=1)} \cdot D_{i=1, j=2} = 1 \cdot 0 = 0$$

$$W_{2,2} = I_{\text{Text}(i=2)} \cdot D_{i=2, j=2} = 0 \cdot 1 = 0$$

$$W_{3,2} = I_{\text{Text}(i=3)} \cdot D_{i=3, j=2} = 0 \cdot 0 = 0$$

$$W_{4,2} = I_{\text{Text}(i=4)} \cdot D_{i=4, j=2} = 0 \cdot 0 = 0$$

(iii) The value of **x**:

$$\sum_{i=1}^4 W_{i,j=3} = W_{1,3} + W_{2,3} + W_{3,3} + W_{4,3} = 0 + 0 + 0 + 0 = 0$$

Where,

$$W_{1,3} = I_{\text{Text}(i=1)} \cdot D_{i=1, j=3} = 1 \cdot 0 = 0$$

$$W_{2,3} = I_{\text{Text}(i=2)} \cdot D_{i=2, j=3} = 0 \cdot 0 = 0$$

$$W_{3,3} = I_{\text{Text}(i=3)} \cdot D_{i=3, j=3} = 0 \cdot 1 = 0$$

$$W_{4,3} = I_{\text{Text}(i=4)} \cdot D_{i=4, j=3} = 0 \cdot 0 = 0$$

(iv) The value of **w**:

$$\sum_{i=1}^4 W_{i,j=4} = W_{1,4} + W_{2,4} + W_{3,4} + W_{4,4} = 0 + 0 + 0 + 0 = 0$$

Where,

$$W_{1,4} = I_{\text{Text}(i=1)} \cdot D_{i=1, j=4} = 1 \cdot 0 = 0$$

$$W_{2,4} = I_{\text{Text}(i=2)} \cdot D_{i=2, j=4} = 0 \cdot 0 = 0$$

$$W_{3,4} = I_{\text{Text}(i=3)} \cdot D_{i=3, j=4} = 0 \cdot 0 = 0$$

$$W_{4,4} = I_{\text{Text}(i=4)} \cdot D_{i=4, j=4} = 0 \cdot 1 = 0$$

This gives the transpose of the binary vector as the output,

z	y	x	w
1	0	0	0

In the same way, output binary vectors of the input given is as follows:

That is $A^I \Rightarrow$

0	0	1	0
1	0	0	0
1	0	0	0
0	0	1	0
1	0	0	0

That is the transpose of inputted Boolean Matrix A of order 4×5 is A^I of order 5×4 .

IV.SUMMARY

This paper presents an illustration oriented study on the algorithm for Boolean Matrix transpose. Here, we observe the algorithm follows parallel computing strategy.

REFERENCES

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