About 2- Isolate Inclusive Sets In Graphs

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Abstract:

In this paper we introduce in new concept call 2-isolate inclusive sets in graphs. Every 2-isolate inclusive set is an isolate inclusive set of G. We characterize maximal 2-isolate inclusive set of a graph. We deduce that every maximal 2-isolate inclusive sets of G is a distance-2 dominating set of G. We also define 2-isolate inclusive number of a graph and we observe that it is less then or equal to isolate inclusive number of the graph. We also prove that if the $\langle S \rangle$ has the maximum number of 2-isolated vertices among all the 2-isolate inclusive sets then S is a maximum 2-packing of G. We also prove several other related results.

Keywords: 2-isolated vertex, 2-packing, 2-isolate inclusive set, maximum 2-isolate inclusive set, maximal 2-isolate inclusive set, distance-k dominating set, distance-2 open neighbourhood, 2-degree of vertex,2-isolate distance-2 dominating set, distance-2 private neighbourhood.

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I. INTRODUCTION

The concept of isolate inclusive set was introduced in [3]. Several interesting results have been proved about isolate inclusive sets. We now introduce a new concept called 2-isolate inclusive sets in graphs. If $S \subset V(G)$ and $v \in S$ then v is said to be 2-isolated vertex in S. If d(v, u) > 2 distance between v and u stractly grather then 2, for all $u \in S$ if $u \neq v$. A set S of vertices is said to be 2-isolate inclusive set if it contains a 2isolated vertex. We consider maximum 2-isolate inclusive sets and maximal 2-isolate inclusive sets in graphs. We prove that every maximal 2-isolate inclusive set is a distance-2 dominating set of G.

We observe that isolate inclusive number [3] of any graph is at least as be as 2-isolate inclusive number of the graph. We further observe that if a graph has an isolated vertex then it has only one 2-isolate inclusive set namely vertex set of the graph.

Here, we also introduce 2-isolate distance-2 dominating set in graphs. We further study the effect of removing a vertex from the graph on 2-isolate distance-2 domination number of a graph. We also consider the operation of edge removal an observe its effect on 2-isolate distance-2 domination number of the graph.

II. PRELIMINARIES AND NOTATIONS

If G is a graph then V(G) denotes the vertex set of the graph G and E(G) denotes the edge set of the graph G. If v is vertex of the graph G then G - v is the subgraph of G induced by all the vertices different from v.

We will consider only simple undirected graphs with finite vertex set.

III. DEFINITIONS AND EXAMPLES

Definition 3.1 (2-isolated vertex) : Let G be a graph and $S \subset V(G)$ a vertex $v \in S$ is said to be 2-isolated vertex of S if d(v, u) > 2, for all $u \in S$ with $u \neq v$.

Definition 3.2 (2-isolate inclusive set) :

Let G be a graph and $S \subset V(G)$ then S is said to be 2-isolate inclusive set if S contains a 2-isolated vertex.

It is obvious that every 2-isolated vertex of S is an isolated vertex of S and every 2-isolate inclusive is an isolate inclusive set.

Example 3.3: Consider the path graph P_5 with 5 vertices $\{1, 2, 3, 4, 5\}$

1 2 3 4 5

Let $S = \{1, 2, 5\}.$

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In S, 5 is a 2-isolated vertex of S and therefore S is a 2-isolate inclusive set.

Consider the path graph P_5 as above. And let $T = \{1, 2, 4\}$. Then 4 is an isolate in T but it is not 2-isolate of T.

Remark 3.4:

Let G be a graph and $v \in V(G)$. Then v is 2-isolated vertex of V(G) if and only if v is an isolated vertex of G. *Definition 3.5 (2-packing)* :[7]

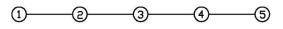
Let G be a graph and $S \subset V(G)$ then S is said to be a 2-packing if d(u, v) > 2, for all $u, v \in S$. Let G be a graph. A 2-packing of G with maximum cardinality is called maximum 2-packing of G. The cardinality of a maximum 2-packing is called the packing number of G and it is denoted as $\delta(G)$. *Remark 3.6:*

Let *S* be a 2-packing of G then every vertex of *S* is a 2-isolated vertex of *S*.

Definition 3.7 (maximum 2-isoinc set) :

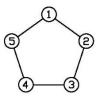
Let G be a graph. A 2-isolate inclusive set with maximum cardinality is called a maximum 2-isolate and its cardinality is denoted as $\beta_{2is}(G)$.

Example 3.8: Consider the path graph P₅ with 5 vertices {1, 2, 3, 4, 5}



Let $S = \{1, 2, 5\}$ then S is a maximum 2-isoinc set and $\beta_{2is}(5) = 3$. $\beta_{2is}(G) = |S| \le \beta_{is}(G)$ $\beta_{2is}(G) \le \beta_{is}(G)$

Example 3.9: Consider the cycle graph C₅ with 5 vertices {1, 2, 3, 4, 5}

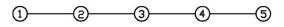


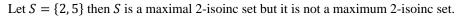
Let $S = \{1, 2, 3, 5\}$ then $\beta_{is}(G) = 4$ and Let $S = \{1, 2, 5\}$ then $\beta_{2is}(G) = 3$. Therefore $\beta_{2is}(G) < \beta_{is}(G)$.

Definition 3.10 (maximal 2-isoinc set) :

Let *G* be a graph and $S \subset V(G)$ be a 2-isoinc set then S is said to be a maximal 2-isoinc set if it is not properly contain in any isoinc set. Obviously every maximum 2-isoinc set is a maximal isoinc set.

Example 3.11: Consider the path graph P_5 with 5 vertices {1, 2, 3, 4, 5}





Definition 3.12 (distance-k dominating set) :[7]

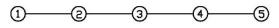
Let G be a graph and $S \subset V(G)$. Then S is said to be a distance-k dominating set, if for every $v \in V(G) - S$, there is a vertex u in S such that $d(v, u) \leq k$, $(k \geq 1)$.

Definition 3.13 (distance-2 open neighbourhood) :[7] Let *G* be a graph and $v \in V(G)$. Then the distance-2 open neighbourhood of v, $N_2(v) = \{u \in V(G) \ni u \neq v \& d(u, v) \le 2\}$ also the distance-2 close neighbourhood of v, $N_2[v] = N_2(v) \cup \{v\}$.

Definition 3.14 (2-degree of vertex) :

Let *G* be a graph and $v \in V(G)$. Then the cardinality of $|N_2(v)|$ will be called the 2-degree of vertex. The minimum 2-degree of a graph *G* will be denoted as $\delta_2(G)$.

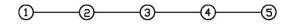
Example 3.15: Consider the path graph P₅ with 5 vertices {1, 2, 3, 4, 5}



Note that if v is an isolated vertex then 2-degree of v = 0. Conversely, also if 2-degree of v = 0 then v is an isolated vertex.

If d(v) = 1 then it is not necessary that $d_2(v) = 1$.

Example 3.16: Consider the path graph P₅ with 5 vertices {1, 2, 3, 4, 5}



Here, d(1) = 1 but $d_2(1) = 1$.

Definition 3.17 (distance-2 dominating set) :[7]

Let G be a graph and $S \subset V(G)$. Then S is said to be a distance-2 dominating set if for every $v \in V(G) - S$, there is a vertex u in S such that $d(v, u) \leq 2$.

A distance-2 dominating set with minimum cardinality is called a minimum distance-2 dominating set.

The cardinality of a minimum distance-2 dominating set is called the distance-2 domination number of the graph and it is denoted as $\gamma_{\leq 2}(G)$.

Definition 3.18 (minimal distance-2 dominating set) :[7]

A distance-2 dominating set *S* is said to be a minimal distance-2 dominating set if $S - \{v\}$ is not a distance-2 dominating set, for each $v \in S$.

Note that every minimum distance-2 dominating set is minimal distance-2 dominating set.

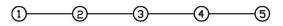
Definition 3.19 (2-isolate distance-2 dominating set) :

Let *G* be a graph and $S \subset V(G)$. Then *S* is said to be a 2-isolate distance-2 dominating set if (1) *S* is a distance-2 dominating set (2) < S > contains a 2-isolated vertex.

Let G be a graph. A 2-isolate distance-2 dominating set with minimum cardinality is called a minimum 2-isolate distance-2 dominating set. It is denoted as $\gamma_{0\leq 2}$ -set.

The cardinality of a $\gamma_{0\leq 2}$ -set is called the 2-isolate distance-2 domination number of the graph and it is denoted as $\gamma_{0\leq 2}(G)$.

Example 3.20: Consider the path graph P_5 with 5 vertices {1, 2, 3, 4, 5}



Let $S = \{2, 5\}$ then S is a minimum 2-isolate distance-2 dominating set of a graph.

Let $T = \{1, 4, 7\}$ then T is a minimum 2-isolate distance-2 dominating set of a graph. Note that |T| < |S|.

Definition 3.21 (distance-2 private neighbourhood) :

Let G be a graph and $S \subset V(G)$ and $v \in S$. Then distance-2 private neighbourhood of v with respect to the set S is equal to $P_{rnd\,2}[v,S] = \{w \in V(G) \ni N_2[w] \cap S = \{v\}\}.$

Remark 3.22:

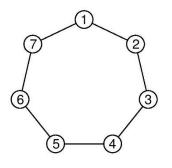
Let G be a graph, $v \in V(G)$ and $v \in S$.

(1) If d(v, u) > 2 for every $u \in S$ with $u \neq v$ then $v \in P_{rnd 2}[v, S]$.

(2) If $x \in S$ and $x \neq v$ then $x \notin P_{rnd 2}[v, S]$.

(3) If $w \in V(G) - S$ then $w \in P_{rnd2}[v, S]$ if and only if v is the only vertex in S whose distance from w is ≤ 2 .

Example 3.23: Consider the cycle graph C₇ with 7 vertices {1, 2, 3, 4, 5, 6, 7}



Let $S = \{2, 5\}$. Let $v = \{2\}$.

- (1) $v \notin P_{rnd\,2}[v, S]$ because $v \in S, 1 \in S$ and $d(v, 1) \leq 2$.
- (2) $1 \notin P_{rnd 2}[v, S]$ because $1 \in S$ and $1 \neq S$.
- (3) $3 \notin S$ but $3 \in P_{rnd 2}[v, S]$ because $d(3, 2) \le 2$ and also $d(3, 1) \le 2$.
- (4) $4 \notin S$ but $4 \in P_{rnd 2}[v, S]$ because d(4, v) = 2 and also d(4, 1) = 3. Which is > 2.
- Similarly $5 \in P_{rnd 2}[v, S]$.
- (5) $6 \notin S$ but $6 \notin P_{rnd 2}[v, S]$ because d(6, v) = 3 > 2.

IV. MAIN RESULT

Proposition 4.1: Let G be a graph, $S \subset V(G)$ and $v \in S$ then v is a 2-isolated vertex of S if and only if $N(v) \cap N(u) = \emptyset$, for every $u \in S$ with $u \neq v$.

Proof: Suppose v is 2 -isolated in S and suppose u is in S, $N(v) \cap N(u) \neq \emptyset$. Let $2 \in N(v) \cap N(u)$. Then v is adjacent to 2 and 2 is adjacent to u. Therefore $d(v, u) \leq d(v, z) + d(z, u) = 1 + 1 = 2$. This is a contradiction. Therefore for every $u \in S$ with $u \neq v$. $N(v) \cap N(u) = \emptyset$.

Conversely, suppose condition is holds. Then d(v, u) > 2, for every $u \in S$ with $u \neq v$. Therefore v is 2 isolated in S

Proposition 4.2: Let S be a 2-isoinc set and $v \in V(G) - S$. Then $S \cup \{v\}$ is not a 2-isoinc set if and only if $d(v, u) \leq 2$, for all 2-isolated verties u of S.

Proof: Suppose $S \cup \{v\}$ is not a 2-isoinc set.

Then $d(u, x) \le 2$, for every isolate *u* of *S* and for some $x \in S \cup \{v\}$ but d(u, w) > 2. For each $w \in S$ therefore $d(u, v) \le 2$, for each 2-isolated vertex *u* of *S*. Conversely, suppose $d(u, v) \le 2$, for each 2-isolated vertex *u* of *S*. Then obviously $S \cup \{v\}$ does not have any 2-isolated vertex.

Theorem 4.3: Let G be a graph and S be a 2-isoinc set of G. Then S is a maximal 2-isoinc set if and only if for every $v \in V(G) - S, S \cup \{v\}$ is not a 2-isoinc set of G.

Proof: Suppose *S* is maximal 2-isoinc set.

Let $v \in V(G) - S$. Since $S \cup \{v\}$ properly contain $S, S \cup \{v\}$ cannot be a 2-isoinc set of G.

Conversely, suppose the condition holds.

Suppose $T \subset V(G)$ is such that S is a proper subset of T if $T = S \cup \{v\}$ for some $u \in V(G) - S$ then by the given condition T cannot be a 2-isoinc set of G.

Therefore we may assume that $|T| - |S| \ge 2$.

Let $v \in T - S$ by the given condition $S \cup \{v\}$ does not have any 2-isolated vertex.

Let $u \in T - S$ be such that $S \cup \{u\}$ does not have any 2-isolated vertex.

Continuing this way, we see that $T = S \cup \{x_1, x_2, \dots, x_k\}$ is not a 2 -isoinc set $T - S = \{x_1, x_2, \dots, x_k\}$.

Thus, the theorem is prove.

Theorem 4.4: Let G be a graph with $\beta_{2is}(G) \ge 2$ then $\beta_{2is}(G) < \beta_{is}(G)$.

Proof: Let S be a maximum 2-isoinc set of G.

Let $v, v' \in S$ and assume that v is 2-isolated vertex of S. Then d(v, v') > 2. Suppose d(v, v') = 3.

Let vu_1u_2v' with a shortest path joining v & v' in *G*. Then $u_1 \notin S$ and $u_2 \notin S$.

Also note that u_2 is not adjacent to v.

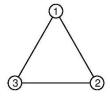
Let $S_1 = S \cup \{u_2\}$. Since v is not adjacent to u_2 and v is also not adjacent to any vertex of S. It follows that v is not adjacent to any vertex of S_1 .

Therefore S_1 is an isoinc set of G.

Therefore $\hat{\beta}_{is}(G) \ge |S_1| > |S| = \beta_{2is}(G)$. Thus $\beta_{2is}(G) < \beta_{is}(G)$.

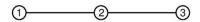
Remark : If $\beta_{2is}(G) = 1$ then the above theorem is not true.

Example 4.5: Consider the triangle with vertices {1, 2, 3}



Then $\beta_{is}(G) = 1$ and $\beta_{2is}(G) = 1$. However, it is also not true that $\beta_{is}(G) = \beta_{2is}(G)$ if $\beta_{2is}(G) = 1$.

Example 4.6: Consider the path graph $G = P_3$ with 3 vertices {1, 2, 3}



Here, $\beta_{2is}(G) = 1$ and $\beta_{is}(G) = 2$.

Theorem 4.7: Let G be a graph and $S \subset V(G)$ be a 2-isoinc set of G. Then S is a maximal 2-isoinc set if and only if for each $v \in V(G) - S$, $d(v, u) \leq 2$ for every 2-isolated vertex u of S.

Proof: Suppose S is a maximal and $v \in V(G) - S$. Let $S_1 = S \cup \{v\}$. Then S_1 does not have any 2-isolated vertex. This means that $d(v, u) \le 2$ for every 2-isolated vertex u of $S, S \cup \{v\}$ cannot have any 2-isolated vertex. Thus, S is a maximal 2-isolate of G.

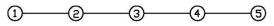
Corollary 4.8: Let *G* be a graph and *v* be an isolated vertex of *G*. If *S* is any maximal 2-isoinc set then $v \in S$.

Proof: Suppose for some maximal 2-isoinc set $S, v \notin S$. Then $d(v, u) \leq 2$, for every 2-isolated vertex u of S. This implies that v is not a isolated vertex in G. Which is a contradiction. Thus the result is proved.

Corollary 4.9: Let G be a graph and $S \subset V(G)$ be a maximal 2-isoinc set of G. Then S is a distance-2 dominating set of G.

Proof: Let $v \in V(G) - S$. By theorem-4.5, $d(v, u) \le 2$ for every 2-isolated vertex u of S. There is a vertex u in S which is a 2-isolated vertex of S. Therefore $d(v, u) \le 2$. Thus S is a distance-2 dominating set of G.

Example 4.10: Consider the path graph P_5 with 5 vertices {1, 2, 3, 4, 5}



Note that if v is an isolated vertex then $\delta_2(v) = 0$. Conversely also if $\delta_2(v) = 0$ then v is an isolated vertex. If d(v) = 1 then it is not necessary that $d_2(v) = 1$.

Theorem 4.11: Let *G* be a graph and $v \in V(G) \ni d_2(v) = \delta_2(v)$. Let $T = V(G) - N_2(v)$ then *T* is a maximum 2-isoinc set of *G*.

Proof: Obviously, *T* is a 2-isoinc set of *G*. Suppose *T* is not a maximum 2-isoinc set of *G*. Then there is a 2-isoinc set of *G* such that |S| > |T|. Let *x* be any 2-isolated vertex of *S*. Then (1) $d_2(v) = \delta_2(v)$. (2) $N_2(x) \subset V(G) - S$. Therefore $S \subset V(G) - N_2(x) \subset V(G) - N_2(x) = T$ and Therefore $|S| \leq |T|$. Which is a contradiction. Therefore *T* is a maximum 2-isoinc set of *G*.

Theorem 4.12: Let G be a graph and T be a maximum 2-isoinc set of G. Then there is $v \in T \ni d_2(v) = \delta_2(G)$ and $T = V(G) - N_2(v)$.

Proof: Let v be any isolated vertex of $\langle T \rangle$ then $N_2(v) \subset V(G) - T$ or $T \subset V(G) - N_2(v)$. Now $V(G) - N_2(v)$ is an isolate of G and also note that $|T| \leq |V(G) - N_2(v)|$. Therefore *T* be a maximum 2-isoinc set of *G*. $|T| = |V(G) - N_2(v)|$. Since $T \subset V(G) - N_2(v)$, $T = V(G) - N_2(v)$. Suppose $d_2(v) > \delta_2(G)$. Let *x* be a vertex of *G* such that $d_2(x) = \delta_2(G)$, by above theorem-4.8 $V(G) - N_2(v)$ is a maximum 2-isoinc set of *G*. Since $d_2(v) > d_2(x)$. $|V(G) - N_2(v)| < |V(G) - N_2(x)|$. Then this implies that $V(G) - N_2(v)$ is not a maximum 2-isoinc set of *G*. Which is a contradiction. Thus $d_2(v) = \delta_2(G)$.

Corollary 4.13: Let G be a graph and v be an isolated vertex then V(G) is the only maximum 2-isoinc set of G.

Proof: Since v is an isolated vertex, $N_2(v) = \emptyset$ and by the above theorem, $V(G) - N_2(v) = V(G)$ is a maximum 2-isoinc set of G. If S is a proper subset of V(G) then obviously S cannot be a maximum 2-isoinc set of G. Thus V(G) is the only maximum 2-isoinc set of G.

Theorem 4.14: Let G be a graph and $S \subset V(G)$ be such that $\langle S \rangle$ has the maximum number of 2-isolated vertex among all the 2-isolate of G. Then S is a maximum 2-paking of G.

Proof: Let S_1 be the set of all 2-isolated vertices of $\langle S \rangle$. Let *M* be a maximum 2-paking of *G*. Then $|S_1| \ge |M| = \delta(G)$. Note that S_1 itself 2-paking of *G* with $|S_1| \ge |M|$. Therefore $|S_1| = |M|$. Thus S_1 is a maximum 2-paking of *G*. Suppose $|S| > |S_1|$. Let $x \in S \ni x \notin S$. Since S_1 is a maximum 2-paking, $d(x, y) \le 2$, for some *y* in S_1 but then this means that *y* is not a 2-isolated vertices in the $\langle S \rangle$. Which is a contradiction. Thus $|S_1| = |S|$ and hence $S_1 = S$ and Therefore *S* is a maximum 2-paking of *G*.

Proposition 4.15: If S is a maximum 2-isoinc set then S is a 2-isolate distance-2 dominating set of G.

Proof: Let $v \in V(G) - S$. Then by theorem-4.5, there is a vertex u in S such that $d(v, u) \le 2$. Thus S is a 2-isolate distance-2 dominating set of G.

Theorem 4.16: Let S is a maximum 2-isoinc set of G (1) For each 2-isolated vertex v of S, $N_2(v) = V(G) - S$. (2) If u and v are 2-isolates of S then $d_2(u) = d_2(v)$.

Proof: (1) Let v be an 2-isolated vertex of S then $N_2(v) \subset V(G) - S$. Let $x \in V(G) - S$. Since S is maximal, $d(x, w) \leq 2$, every 2-isolated vertex w of S. Therefore $x \in N_2(v)$. Thus $N_2(v) = V(G) - S$.

(2) If *u* and *v* be two 2-isolates of *S* then $d_2(u) = |N_2(u)| = |V(G) - S| = |N_2(v)| = d_2(v).$ Thus $d_2(u) = d_2(v).$ **Remark 4.17:** Let G be a graph and v be an isolated vertex in G. It is obvious that d(x, y) in G is equal to d(x, y) in G - v as v is an isolated vertex in G.

Theorem 4.18: Let *G* be a graph and *v* be an isolated vertex in *G*. Then $\gamma_{\leq 2}(G) \leq \gamma_{\leq 2}(G-v)$ if and only if for every minimum 2-isolate distance-2 dominating set *S* of *G*. The following condition is satisfied **C**: *v* is the only 2-isolated vertex in the $\langle S \rangle$.

Proof: Suppose $\gamma_{\leq 2}(G) \leq \gamma_{\leq 2}(G - v)$. Let *S* be any minimum 2-isolate distance-2 dominating set of *G*. Since *v* is an isolated vertex of *G*, $v \in S$. Suppose there is vertex $v' \in S \ni v' = v \& v'$ is 2-isolated vertex of *S*. Now let $S_1 = S - \{v\}$. Consider the subgraph G - v. Let $x \notin S_1$. Then there is a vertex *y* in *S* such that $d(x, y) \leq 2$ in *G*. Obviously $y \neq v$. Therefore $d(x, y) \leq 2$ in G - v also. Thus S_1 is a 2-isolate distance-2 dominating set of G - v. Therefore $\gamma_{\leq 2}(G - v) \leq |S_1| < |S| = \gamma_{\leq 2}(G)$. This contradict the hypothesis that $\gamma_{\leq 2}(G) \leq \gamma_{\leq 2}(G - v)$. Therefore *v* is the only 2-isolated vertex of *S*.

Conversely, suppose the condition is satisfied.

For any minimum 2-isolate distance-2 dominating set of G.

Let *T* be any set of vertices of G - v such that $|T| < \gamma_{\leq 2}(G)$.

Suppose *T* is a 2-isolate distance-2 dominating set in G - v.

Let x be any vertex of G such that $x \neq v$ and $x \notin T$. There is a vertex y in T such that $d(x, y) \leq 2$ in G - v. Since v is an isolated vertex, by the above remark $d(x, y) \leq 2$ in G also.

Note let $S = T \cup \{v\}$ then *S* is a minimum 2-isolate distance-2 dominating set of *G*.

Then $|S| = \gamma_{\leq 2}(G)$ and *S* contains two 2-isolated vertices including *v*.

Which is a contradicyion.

Therefore if $|T| < \gamma_{<2}(G)$ then *T* cannot be 2-isolate distance-2 dominating set of *G*.

Therefore $|T| \ge \gamma_{\le 2}(G)$.

Therefore $\gamma_{\leq 2}(G - v) \geq \gamma_{\leq 2}(G)$.

Theorem 4.19: Let *G* be a graph and *v* be an isolated vertex of *G*. Then $\gamma_{\leq 2}(G - v) < \gamma_{\leq 2}(G)$ if and only if there is a minimum 2-isolate distance-2 dominating set *S* of *G* such that *S* contains an isolate different from *v*.

Proof: Suppose $\gamma_{\leq 2}(G - v) < \gamma_{\leq 2}(G)$.

Let S_1 be a minimum 2-isolate distance-2 dominating set of G - v.

Then S_1 cannot be 2-isolate distance-2 dominating set of *G* because $|S_1| = \gamma_{\leq 2}(G - v) < \gamma_{\leq 2}(G)$. Let $S = S_1 \cup \{v\}$.

Let x be a vertex of S_1 such that $x \notin S$ and $x \notin S_1$ also. Since S_1 is a 2-isolate distance-2 dominating set of G - v, $d(x, y) \le 2$ in G - v, for some y in S. Then $d(x, y) \le 2$ in G also.

Thus *S* is a 2-isolate distance-2 dominating set of *G* such that $v \in S$. Since $|S| = |S_1| + 1$, *S* is a minimum 2-isolate distance-2 dominating set of *G*.

Let v' be if 2-isolated vertex of S_1 then v' is also 2-isolated vertex of S as v is an isolated vertex of G.

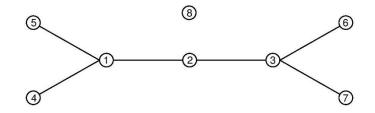
Thus S is a minimum 2-isolate distance-2 dominating set of G which contains an isolate different from v.

Conversely, suppose there is a minimum 2-isolate distance-2 dominating set *S* of *G* such that *S* contains a 2-isolate different from *v*. Since *v* is an isolated vertex in $G, v \in S$.

Let $S_1 = S - \{v\}$. Then S_1 contains a 2-isolate (which is different from v).

Therefore S_1 is a 2-isolate distance-2 dominating set of G - v.

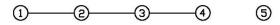
Thus $\gamma_{\leq 2}(G - v) \leq |S_1| < |S| = \gamma_{\leq 2}(G)$. Thus $\gamma_{\leq 2}(G - v) < \gamma_{\leq 2}(G)$. *Example 4.20:* Consider the graph *G* with vertices {1, 2, 3, 4, 5, 6, 7, 8} mansion blow.



We may note that the set $S = \{1, 2, 8\}$ is a minimum 2-isolate distance-2 dominating set of *G*. Also note that for any minimum set *T* of *G*, 8 is the only 2-isolated vertex of *T*.

Now consider the subgraph G - 8. Consider the set $S_1 = \{1, 6, 7\}$. Then S_1 is a minimum 2-isolate distance-2 dominating set of G - 8. Thus $\gamma_{\leq 2}(G - 8) = \gamma_{\leq 2}(G)$.

Example 4.21: Consider the graph *G* with vertices {1, 2, 3, 4, 5} mansion blow.



Let $S = \{1, 4, 5\}$. Then S is a minimum 2-isolate distance-2 dominating set of G. Note that S contains 2-isolate different from 5.(infact 1 and 4 are both 2-isolates of S)

Now consider the subgraph G - 5. Let $T = \{1, 4\}$. Then T is a minimum 2-isolate distance-2 dominating set of G - 5. Therefore thus $\gamma_{\leq 2}(G - 5) = 2 < 3 = \gamma_{\leq 2}(G)$.

Corollary 4.22: Let *G* be a graph and v_1, v_2, \dots, v_k be all the isolated vertices of $G(k \ge 2)$. Then $\gamma_{\le 2}(G - v_i) < \gamma_{\le 2}(G)$, for all $i = 1, 2, \dots, k$.

Proof: Let *S* be a minimum 2-isolate distance-2 dominating set of *G*. Then $v_i \in S$, for every i = 1, 2, ..., k. Then by above corollary, $\gamma_{\leq 2}(G - v_i) < \gamma_{\leq 2}(G)$, for all i = 1, 2, ..., k. ▮

Corollary 4.23: If there is a 2-isolated vertex v such that $\gamma_{\leq 2}(G - v) \geq \gamma_{\leq 2}(G)$ then the graph has only one 2-isolated vertex namely v.

Proof: Obvious .

Theorem 4.24: Let *G* be a graph and *v* be a non isolated vertex in *G*. Then $\gamma_{\leq 2}(G - v) > \gamma_{\leq 2}(G)$ if and only if the following two conditions are satisfied.

(1) For every a minimum 2-isolate distance-2 dominating set *S* of *G*, $d(v, S) \le 1$

(2) There is no subset S of $V(G) - N_2(v)$ such that $|S| \le \gamma_{\le 2}(G)$ and S is a minimum 2-isolate distance-2 dominating set of G - v.

Proof: Suppose $\gamma_{\leq 2}(G - v) > \gamma_{\leq 2}(G)$. (1) Let *S* be a minimum 2-isolate distance-2 dominating set of *G*. If $v \in S$ then d(v, S) = 0 < 1. Suppose $v \notin S$. Consider the subgraph G - v. Since $|S| \leq \gamma_{\leq 2}(G - v)$, *S* cannot be 2-isolate distance-2 dominating set of G - v. Note that any 2-isolate of *S* in *G* is also 2-isolate of *S* in G - v. Therefore the $\langle S \rangle$ in G - v contains 2-isolated vertices. It follows that *S* is not a distance-2 dominating set of G - v. Therefore there is a vertex *x* in G - v such that $x \notin S$ and $d(x, S) \ge 3$ in G - v. Since *S* is a distance-2 dominating set of *G*, $d(x, S) \ge 2$ in *G*. Let *P* be a path joining *x* to some vertex *z* of *S* such that length of $P \le 2$. Obviously this path contain *v*.Since $x \ne v, x$ is adjacent to *v* and *v* is adjacent to *z* in *G*.

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Thus d(v, S) = 1.
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Therefore from both the above cases it follows that $d(v, S) \leq 1$.

(2) Suppose there is a set $S \subset V(G) - N_2[v]$ such that $|S| \leq \gamma_{\leq 2}(G)$ and *S* is a 2-isolate distance-2 dominating set of G - v. Then $\gamma_{\leq 2}(G - v) \leq |S| \leq \gamma_{\leq 2}(G)$. Which implies that $\gamma_{\leq 2}(G - v) \leq \gamma_{\leq 2}(G)$. which is a contradiction. Therefore (2) is also proved.

Thus the theorem is proved.

Now we state and prove a necessary and sufficient condition under which 2-isolate distance-2 domination number decreases when vertex is remove from the graph.

Theorem 4.25: Let G be a graph and $v \in V(G)$. Then $\gamma_{0 \leq 2}(G - v) > \gamma_{0 \leq 2}(G)$ if and only if there is a minimum 2-isolate distance-2 dominating set of S such that

(1) *S* contains a 2-isolate different from *v*. (2) $v \in S$ and $P_{rnd 2}[v, S] = \{v\}$.

Proof: Suppose $\gamma_{0 \leq 2}(G - v) < \gamma_{0 \leq 2}(G)$.

Let S_1 be a minimum 2-isolate distance-2 dominating set of G - v. Let z be a 2-isolate of S_1 in G - v. Since $|S_1| = \gamma_{0 \leq 2}(G - v) < \gamma_{0 \leq 2}(G)$, S_1 cannot be a 2-isolate distance-2 dominating set of G.

Therefore there is a vertex x in G such that $d(x, S_1) > 2$ in G.

If $x \neq v$ then $d(x, S_1) \leq 2$ in G - v because S_1 is a 2-isolate distance-2 dominating set of G - v. Then $d(x, S_1)$ in G is ≤ 2 .

Which is a contradiction.

Therefore $x \neq v$ is not possible.

Therefore x = v and $d(x, S_1) > 2$ in *G*.

Let $S = S_1 \cup \{v\}$. Then $v \in S$. Since d(v, z) > 2 in G, z is also a 2-isolate of S in G.

If $y \in V(G)$ and $y \notin S$ then as prove (1) above $d(y, S_1) \le 2$ in *G*.

Therefore $d(y,S) \le 2$ in *G*. Therefore *S* is a 2-isolate distance-2 dominating set of *G* containing *v*. Since $d(v,S_1) > 2$, d(v,u) > 2, for all $u \in S$ with $u \ne v$.

Therefore $v \in P_{rnd 2}[v, S]$.

Let $T \in V(G) - S$ such that $d(T, v) \leq 2$. Now $T \notin S_1$ and S_1 is a distance-2 dominating set of G - v. Therefore there is a vertex T' in S_1 such that $d(T, T') \leq 2$ in G also.

Thus we have proved that $d(T, v) \leq 2$ in *G*.

Therefore $T \notin P_{rnd 2}[v, S]$.

Thus $P_{rnd 2}[v, S] = \{v\}$. Also note that *S* contain a 2-isolate different from *v*.

Conversely, suppose there is a minimum 2-isolate distance-2 dominating set S of G such that

(1) *S* contain a 2-isolate different from v.

(2) $P_{rnd 2}[v, S] = \{v\}.$

Let $S_1 = S \cup \{v\}$. Let z be 2-isolate of S different from v. Then $z \in S_1$ and z is a 2-isolate of S_1 in G - v. Let x be any vertex of G - v such that $x \notin S$ also. Since S is a distance-2 dominating set of G. There is some y in S such that $d(x, y) \le 2$ in G.

Case(1): suppose y = v.

Since $x \notin P_{rnd 2}[v, S]$. There is a vertex y' in S such that $y' \neq v$ and $d(x, y') \leq 2$ in G. Any path in G joining x to y' whose length is ≤ 2 cannot contain v as an internal vertex because d(x, y') > 2. Therefore Any path joining x to y' in G having length ≤ 2 is also a path in G - v. Therefore $d(x, y') \leq 2$ in G - v.

Case(2): suppose $y \neq v$.

Then $y \in S_1$ and $d(x, y) \le 2$ in *G*. By the same argument given above $d(x, y) \le 2$ in G - v also. Thus we have proved that any vertex of G - v which is not in S_1 satisfies $d(x, y) \le 2$ in G - v, for some y in S_1 .

Therefore S_1 is a 2-isolate distance-2 dominating set of G - v. Therefore $\gamma_{0 \leq 2}(G - v) \leq |S_1| < |S| = \gamma_{0 \leq 2}(G)$. Hence, $\gamma_{0 \leq 2}(G - v) < \gamma_{0 \leq 2}(G)$. Thus the theorem is prove.

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Example 4.26: Consider the path graph P_6 with 6 vertices {1, 2, 3, 4, 5, 6}



Let $S = \{3, 6\}$. It is obvious that *S* is a minimum 2-isolate distance-2 dominating set of *G*. Now consider the graph G - 6 which is the path graph P_5 with vertices $\{1, 2, 3, 4, 5\}$ 2-isolate distance-2 domination number =1 Thus $\gamma_{0 \leq 2}(G - v) < \gamma_{0 \leq 2}(G)$. Observe that $6 \in S$ and $P_{rnd \geq 2}[6, S] = \{6\}$.

Also *S* contains a 2-isolate different from 6.

Corollary 4.27: Let *G* be a graph without isolated vertices. Suppose $\gamma_{0\leq 2}(G-\nu) < \gamma_{0\leq 2}(G)$. Then there is a minimum 2-isolate distance-2 dominating set *S* of *G* such that $\nu \notin S$.

Proof: Since $\gamma_{0 \leq 2}(G - v) < \gamma_{0 \leq 2}(G)$.

There is a minimum 2-isolate distance-2 dominating set of S_1 of G such that (which contains isolate different from v) $P_{rnd 2}[v, S] = \{v\}$.

Since v is not an isolated vertex of G. There is a vertex v' such that $d(v, v') \le 2$. Obviously $v' \notin S_1$. Let $S = (S_1 - \{v\}) \cup \{v'\}$. Then $|S| = |S_1|$.

Let *x* be any vertex of *G* then $x \notin S$.

If x = v then $d(v, v') \le 2$ and $v' \in S$.

If $x \neq v$ then $x \notin S_1$. Since S_1 is a 2-isolate distance-2 dominating set of *G*. There is a vertex z in S_1 such that $d(x, z) \leq 2$.

If z = v then there is another vertex w in S_1 such that $d(x, w) \le 2$ in G because $x \notin P_{rnd 2}[v, S_1]$ then $w \in S$ and $d(x, w) \le 2$.

If $d(x, v') \leq 2$ in G then $v' \in S$ and $d(x, v') \leq 2$.

Thus for any x not in S. There is some vertex y in S such that $d(x, y) \le 2$ in G. This proves that S is a minimum 2-isolate distance-2 dominating set of G. Note that $v \notin S$.

Proposition 4.28: Let *G* be a graph and *v* be an isolated vertex of *G*. If *S* is a 2-isolate distance-2 dominating set of *G* then $v \in S$.

Proof: Suppose $v \notin S$. Then $d(v, S) \leq 2$. Then there is a vertex $u \in S$ such that $d(v, u) \leq 2$. This implies that v is not isolated vertex.

Which is a contradiction.

Therefore $v \in S$.

Theorem 4.29: Let G be a graph and v_1, v_2, \dots, v_k be all the isolated vertices of $G(k \ge 2)$. Then $\gamma_{0\le 2}(G - v_i) < \gamma_{0\le 2}(G)$ for $i = 1, 2, \dots, k$.

Proof: Let *S* be a minimum 2-isolate distance-2 dominating set of *G*. By the above proposition 28, $v_i \in S$, for all i = 1, 2, ..., k. Consider v_i . Now $v_1 \in S$ and *S* also contains a 2-isolate vertex of *S* different from v_1 . Also $P_{rnd 2}[v_1, S] = \{v_1\}$. Therefore $\gamma_{0 \leq 2}(G - v_1) < \gamma_{0 \leq 2}(G)$. Similarly, it can be proved that $\gamma_{0 \leq 2}(G - v_i) < \gamma_{0 \leq 2}(G)$ for i = 1, 2, ..., k. Thus the theorem is proved.

Corollary 4.30: Let G be a graph and v be an isolated vertex of G. If $\gamma_{0\leq 2}(G-v) \geq \gamma_{0\leq 2}(G)$. Then v is the only isolated vertex of G.

Proof: Suppose there is vertex v' of G such that $v' \neq v$ and v' is also isolated vertex. Let S be a minimum 2-isolate distance-2 dominating set of G. Then $v, v' \in S$. Also $P_{rnd 2}[v, S] = \{v\}$. And therefore $\gamma_{0 \leq 2}(G - v) < \gamma_{0 \leq 2}(G)$. Which is a contradiction. Then v is the only isolated vertex of G.

Then *v* is the only isolated vertex of *u*.

Now we consider the operation of edge removal in graph. Proposition 4.31: Let *G* be a graph and *e* be an edge of *G*. If $u, v \in V(G)$. Then d(u, v) in $G - e \ge d(u, v)$ in *G*.

Proof: If there is no path joining u and v in G. Then there is no path joining u and v in G - e, $d(u, v) = \infty$ in G. In this case there is no path joining u and v in G - e also. And therefore $d(u, v) = \infty$ in G - e also. Thus the result is prove in this case.

Suppose d(u, v) = k in G - e, for some positive integer k. Then there is a path of length k joining u and v in G - e. This is also a path joining u and v in G - e.

Therefore d(u, v) in G is \leq the length of the path P which is = k which is = d(u, v) in G - e.

Therefore d(u, v) in $G \le d(u, v)$ in G - e.

Thus the result is proved.

Now we prove the following theorem.

Theorem 4.32: Let *G* be a graph and $e = \{uv\}$ be an edge of *G*. Then $\beta_{2is}(G - e) \ge \beta_{2is}(G)$.

Proof: Let *S* be a maximum 2-isoinc set of *G*. Let $u \in S$ be a 2-isolated vertex of *S*. Then d(u, x) > 2 in *G*, for every $x \in S$ with $x \neq u$. Let *w* be 2-isolated vertex of *S*. Then d(w, x) > 2 in *G*, for every $x \in S$ with $x \neq w$. Then d(w, x) > 2 in G - e also, for every $x \in S$ with $x \neq w$.

Thus w is a 2-isolated vertex of S in G - e.

Moreover, if $a, b \in S$ then d(a, b) > 2 in G, for all $a, b \in S$. Therefore d(a, b) > 2, for all $a, b \in S$ in G - e also.

Thus *S* is a 2-isoinc set in G - e also. Therefore $\beta_{2is}(G - e) \ge |S| = \beta_{2is}(G)$. Therefore $\beta_{2is}(G - e) \ge \beta_{2is}(G)$.

Now we state and prove a necessary and sufficient condition under which 2-isoinc number of a graph increases when an edge is remove on the graph.

Theorem 4.33: Let G be a graph and and $e = \{uv\}$ be an edge of G. Then $\beta_{2is}(G - e) > \beta_{2is}(G)$ if and only if for every maximum 2-isoinc S of G - e. The following conditions are satisfied.

(1) If $u, v \in S$ then for every 2-isolate z of S. There is a vertex w in V(G) - S, which is adjacent to z and u or there is a vertex w' in V(G) - S, which is adjacent to z and v.

(2) If $u \notin S$ and $v \in S$ then for ervery 2-isolate z of S, z is adjacent to u.

(3) If $v \notin S$ and $u \in S$ then for ervery 2-isolate z of S, z is adjacent to v.

Proof: First suppose that $\beta_{2is}(G - e) \ge \beta_{2is}(G)$. Let *S* be a maximum 2-isoinc of G - e. Since $|S| > \beta_{2is}(G)$, *S* cannot be a 2-isoinc of *G*. Therefore if *z* is any 2-isolate of *S* in G - e then *z* cannot be a 2-isolated vertex *S* in *G*. Therefore $d(z, x) \le 2$ in *G* some *x* in *S*. *Case(1):* $u \in S \& v \in S$ It follows that u = x or v = x. If x = u then there is a vertex *w* in V(G) - S such that *w* is adjacent to both u & z. If x = v then there is a vertex *w* in V(G) - S such that *w'* is adjacent to both v & z.

Case(2): $u \notin S \& v \in S$ In this case it follows that x = v. Since d(z, v) > 2 in G - e and $d(z, v) \le 2, z$ must be adjacent to u.

Case(3): $v \notin S \& u \in S$

In this case it follows that x = u. Since d(z, u) > 2 in G - e and $d(z, u) \le 2, z$ must be adjacent to v.

Conversely, suppose conditions (1), (2) and (3) are any one is satisfied, for any maximum 2-isoinc set S of G - e.

Let S be subset of V(G) such that $|S| \ge \beta_{2is}(G - e)$. Suppose S is a 2-isoinc of G, S must be a 2-isoinc set S of G - e also.

Thus here $|S| \ge \beta_{2is}(G - e)$ and S is a 2-isoinc set S of G - e.

This is a contradiction.

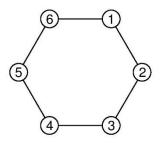
Thus *S* cannot be a 2-isoinc set *S* of *G* if $|S| \ge \beta_{2is}(G - e)$.

Suppose $|S| = \beta_{2is}(G - e)$. Suppose S is a 2-isoinc of G. Now S is also a maximum 2-isoinc of G - e. By the assumption conditions (1), (2) or (3) are satisfied by S and therefore S cannot be a 2-isoinc set of G. Which is again contradiction.

Thus we have prove that $|S| \ge \beta_{2is}(G - e)$ then S cannot be a 2-isoinc set of G.

Therefore if *T* is any maximum 2-isoinc set of *G* then $|T| < \beta_{2is}(G)$. Therefore $\beta_{2is}(G - e) > \beta_{2is}(G)$.

Example 4.34: Let G be the cycle graph C_6 with 6 vertices {1, 2, 3, 4, 5, 6}



In this graph $S = \{1, 4\}, \{2, 5\}, \{3, 6\}$ are the only maximum 2-isoinc sets of *G*. Now consider the graph G - e. Which is the path graph P_6 with 6 vertices $\{1, 2, 3, 4, 5, 6\}$.

In this graph $S_1 = \{1, 4, 5, 6\}$ and $S_1 = \{1, 2, 3, 6\}$ are the only maximum 2-isoinc sets of G - e. Note that $\{1, 6\} \in S_1$ and $\{1, 6\} \in S_2$ also. Let w = 2 then w is the adjacent to 1 for the set S_1 . Let w = 5 then w is the adjacent to 6 for the set S_2 .

Example 4.35: Consider the complete graph K_n with $n \ge 3$.

Then $\beta_{2is}(K_n) = 1$. Remove any edge from the graph K_n then $\beta_{2is}(K_n - e) = 1$. Here, 2-isoinc number does not increases when any edge remove from the graph.

V. CONCLUDING REMARKS

In this paper we have consider 2-isolate inclusive set. It may be possible to study those sets which do not contain 2-isolated vertices. These sets can be studied and can be compared with totally dominating sets.

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