

# Laceability in the Interleaver Graph of Brick Product Graph $C(2n, 1, n)$

R. A. Daisy Singh <sup>1</sup>, Leena N Shenoy <sup>2</sup>

<sup>1</sup> Department of Mathematics, St. Joseph's College Autonomous, Bengaluru, Karnataka, India.

<sup>2</sup> Department of Mathematics, B. N. M. Institute of Technology, Bengaluru, Karnataka, India.

**Abstract** - Interleavers are used as a tool in constructing good turbo codes which are the class of error correcting codes introduced by Berrou and Glavieux in 1993. Interleavers have been used in communication systems such as 3G/4G mobile communications and in satellite communications etc. The interleaver's function is to permute low weight code words in one encoder into high weight code words for the other encoder. In this paper we present the construction of interleaver graphs  $IG_N$ ,  $N \geq 6$  from the brick product graph  $C(2n, 1, n)$ ,  $n \geq 3$  and discuss its Hamiltonian laceability properties.

**Keywords:** Interleaver graphs, Brick product graphs, Hamiltonian connected, Hamiltonian laceability.

## 1 Introduction

Graphs are important structures by using which many applications of the real world such as communication networks and social networks are described. In [1] Arya Mazumdar et.al have proposed a new construction of interleavers for turbo codes from 3-regular hamiltonian graphs. A good interleaver for turbo codes can be constructed from 3-regular Hamiltonian graphs having large girth. The girth of a graph is the shortest cycle contained in a graph. In topological point of view it is important for any interconnection network to have various graph theoretic properties which includes girth.

Using the concept of brick product Brain Alspach et.al. showed in [2] that all cubic Cayley graphs over dihedral groups are Hamiltonian. Also in [3] Brain Alspach et.al. showed that most  $C(2n, 1, r)$  are Hamiltonian laceable for  $3 \leq r \leq 9$ . Inspired by their work we developed the construction of interleaver graph  $IG_N$  for  $N \geq 6$  from the brick product graph  $C(2n, 1, n)$  and explore its Hamiltonian  $t^*$  laceable properties. Further fault tolerance is an important property in the network performance which enables systems to continue to operate properly in the case of failure of some of the components. Sun-Yuan Hsieh et.al.in [4] proposed the edge fault tolerant Hamiltonicity property in the faulty network. Extending this we also explore the edge fault tolerant Hamiltonian  $t^*$  laceability of interleaver graphs  $IG_N$ ,  $N \geq 6$  for all even  $N$  such that  $1 \leq t \leq \text{diam}(G)$ .

## 2 Interleavers from the brick product graph $C(2n, 1, n)$

Let  $G = C(2n, 1, n)$ . Let  $2n = N$ ,  $n \geq 3$  be the number of vertices labelled from 1 to  $N$  placed on a cycle. Initially the fixed cycle in  $G$  becomes the Hamiltonian cycle of the graph.

To construct  $IG_N$  from  $G$  we first draw two chains of  $N$  vertices, one in the upper part and the other in the lower part of the graph. Note that the chain in the upper and lower parts are of length  $N$ . Let  $i$  be the vertex in the chain of upper part and  $j$  be the vertex in the chain of lower part then  $i$  and  $j$  should be connected in  $IG_N$  if there is an edge which connects  $i$  and  $j$  in  $G$  and is not included in the Hamiltonian cycle of  $G$ .

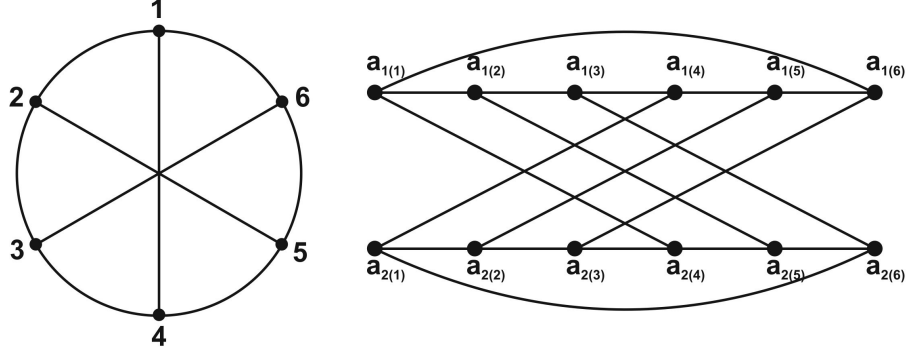


Figure 1 : Construction of  $IG_6$  from  $C(6, 1, 3)$

### 3 Terminologies

We need to first introduce the following terminologies to establish results.

For each vertex  $a_{1i}$  and  $a_{2i}$  of  $IG_N$  where  $1 \leq i < N$ , we shall write:

$$\begin{aligned} a_{1i}\phi[N] &= a_{1i}(a_{1(i+1)})(a_{1(i+2)})(a_{1(i+3)})(a_{1(i+4)})\dots\dots(a_{1(i+N-1)}) \\ a_{1i}\phi^{-1}[N] &= a_{1i}(a_{1(i-1)})(a_{1(i-2)})(a_{1(i-3)})(a_{1(i-4)})\dots\dots(a_{1(i-N+1)}) \\ a_{2i}\Gamma[N] &= a_{2i}(a_{2(i+1)})(a_{2(i+2)})(a_{2(i+3)})(a_{2(i+4)})\dots\dots(a_{2(i+N-1)}) \\ a_{2i}\Gamma^{-1}[N] &= a_{2i}(a_{2(i-1)})(a_{2(i-2)})(a_{2(i-3)})(a_{2(i-4)})\dots\dots(a_{2(i-N+1)}) \end{aligned}$$

Note that the symbols  $\phi[N]$ ,  $\phi^{-1}[N]$ ,  $\Gamma[N]$ ,  $\Gamma^{-1}[N]$  are the paths of order  $N$  where as the symbols  $R[x]$ ,  $R^{-1}[x]$ ,  $S[x]$ ,  $S^{-1}[x]$ ,  $T[x]$ ,  $T^{-1}[x]$ ,  $H[x]$ ,  $H^{-1}[x]$ ,  $L[x]$  are self explanatory.

For  $x = 1$

$$\begin{aligned} R[x] &= a_{11} \rightarrow a_{1N} \\ R^{-1}[x] &= a_{1N} \rightarrow a_{11} \\ S[x] &= a_{21} \rightarrow a_{2N} \\ S^{-1}[x] &= a_{2N} \rightarrow a_{21} \\ L[x] &= a_{12} \rightarrow a_{22} \end{aligned}$$

For  $1 \leq i \leq \frac{N}{2}$ ,  $k = \frac{N}{2} + i$ ;  $x = 1$

$$\begin{aligned} T[x] &= a_{1i} \rightarrow a_{2k} \\ T^{-1}[x] &= a_{2k} \rightarrow a_{1i} \\ H[x] &= a_{2i} \rightarrow a_{1k} \\ H^{-1}[x] &= a_{1k} \rightarrow a_{2i} \end{aligned}$$

**3.1 Definition:** A connected graph  $G$  is said to be Hamiltonian  $t$  laceable (Hamiltonian  $t^*$  laceable) if there exists a Hamiltonian path between every pair (at least one pair) of distinct vertices  $u$  and  $v$  in  $G$  with the property  $d(u, v) = t$  such that  $1 \leq t \leq \text{diam}(G)$  where  $\text{diam}(G)$  is the diameter of  $G$ .

**3.2 Definition:** A connected graph  $G$  is termed as Hamiltonian  $t_o^*$  ( $t_e^*$ ) laceable if there exists in  $G$ , a Hamiltonian path between atleast one pair of its vertices  $u$  and  $v$  with the property  $d(u, v) = t$  for odd (even)  $t$  such that  $1 \leq t \leq \text{diam}(G)$ .

**3.3 Definition:** The graph  $G'$  is  $k$ -edge fault tolerant with respect to a graph  $G$  if every graph obtained by removing any  $k$ -edges from  $G'$  contains  $G$ . Further a graph  $G'$  is  $k$ -edge fault tolerant Hamiltonian laceable if  $G' - F$  remains Hamiltonian laceable for every  $F \subseteq E(G')$  with  $|F| \leq k$ .

**3.4 Definition:** Let  $P$  be a path between the vertices  $a_{v_i}$  and  $a_{v_j}$  in a graph  $G$  and  $P'$  be the path between the vertices  $a_{v_j}$  and  $a_{v_k}$ . Then the path  $P \cup P'$  is obtained by extending  $P$  from  $a_{v_i}$  to  $a_{v_j}$  and then from  $a_{v_j}$  to  $a_{v_k}$  through the common vertex  $a_{v_j}$ . That is if  $P : a_{v_i} \rightarrow a_{v_j}$  and  $P' : a_{v_j} \rightarrow a_{v_k}$  then  $P \cup P' : a_{v_i} \rightarrow a_{v_j} \rightarrow a_{v_k}$ .

## 4 Results

### 4.1 Theorem

Let  $IG_N$  be the interleaver graph constructed from  $C(2n, 1, n)$  for  $n \geq 3$ . Then  $IG_N$ ,  $N \geq 6$  is Hamiltonian  $t_o^*$  laceable for  $1 \leq t \leq \text{diam}(G)$ .

**Proof:** Let the vertex set of  $IG_N$  be  $V = V_1 \cup V_2$  such that  $V_1 = \{a_{i(j)} : i = 1, 1 \leq j \leq N\}$ ,  $V_2 = \{a_{i(j)} : i = 2, 1 \leq j \leq N\}$  and the edge set of  $IG_N$  be  $E = \{b_k : 1 \leq k \leq 3N\}$ . Clearly the diameter of  $IG_N$  is  $D = \frac{N}{2} + 1$ .

#### Case (i)

In  $IG_N$ ,  $N = 4h + 2$  for  $h \geq 1$ , we have  $d(a_{1(1)}, a_{1(2)}) = 1$  and the path  $P : \{a_{1(1)} - T[x] - \Gamma^{-1}[2] - H[x] - \phi^{-1}[2] - H^{-1}[x] - \Gamma^{-1}[2] - H[x] - \phi^{-1}[2] - \{H^{-1}[x] - \Gamma^{-1}[2] - H[x] - \phi^{-1}[2]\}_{h-1} - T[x] - \Gamma^{-1}[2] - T^{-1}[x] - \{\phi^{-1}[2] - T[x] - \Gamma^{-1}[2] - T^{-1}[x]\}_{h-1} - a_{1(2)}\}$  is a Hamiltonian path between the vertices  $a_{1(1)}$  and  $a_{1(2)}$ .

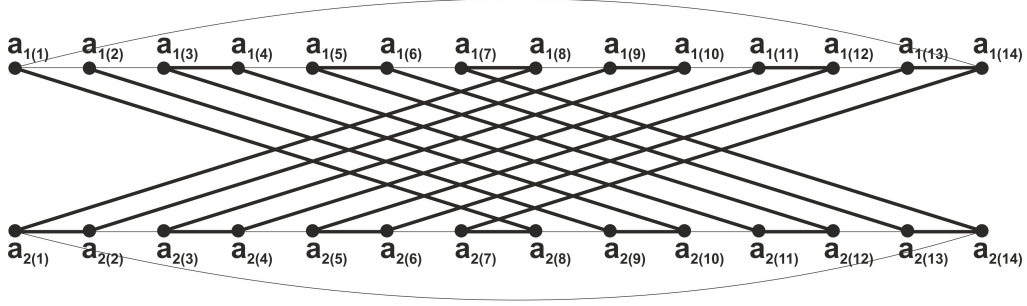


Figure 2 : Hamiltonian path between  $a_{1(1)}$  and  $a_{1(14)}$   $IG_{14}$ .

### Case (ii)

In  $IG_N$ ,  $N = 4h + 2$  for  $h \geq 1$ , we have  $d(a_{1(1)}, a_{1(2h+2)}) = t$  such that  $3 \leq t \leq 2h + 1$  for odd  $t$  and the path  $P : \{a_{1(1)} - \phi[2] - T[x] - \Gamma^{-1}[3] - \{H[x] - \phi^{-1}[2] - H^{-1}[x] - \Gamma^{-1}[2]\}_h - S[x] - \Gamma^{-1}[2h - 1] - T^{-1}[x] - \phi[2h] - a_{1(2h+2)}\}$  is a Hamiltonian path between the vertices between  $a_{1(1)}$  and  $a_{1(2h+2)}$ .

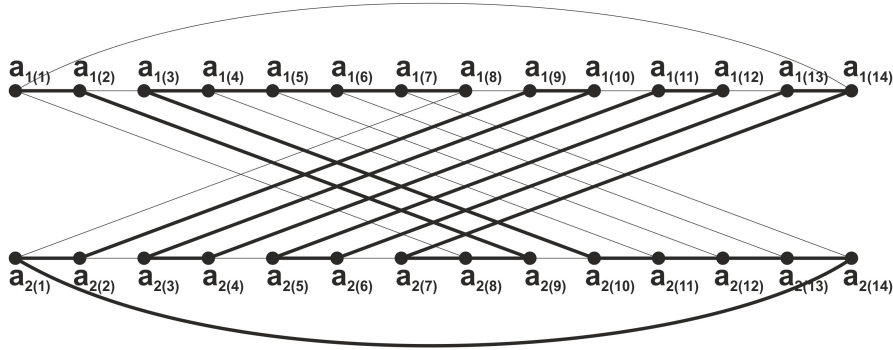


Figure 3 : Hamiltonian path between  $a_{1(1)}$  and  $a_{1(8)}$   $IG_{14}$ .

### Case (iii)

In  $IG_N$ ,  $N = 4h + 6$  for  $h \geq 1$ , we have  $d(a_{1(1)}, a_{1(2h+2)}) = t$  such that  $3 \leq t \leq 2h + 1$  for odd  $t$  and the path  $P : \{a_{1(1)} - \phi[2] - T[x] - \Gamma^{-1}[3] - \{H[x] - \phi^{-1}[2] - H^{-1}[x] - \Gamma^{-1}[2]\}_{h+s} - H[x] - \phi^{-1}[2] - T[x] - \{\Gamma^{-1}[2] - T^{-1}[x] - \phi^{-1}[2] - T[x]\}_{h-1} - \Gamma[2h + 1] - T^{-1}[x] - \phi[2h] - a_{1(2h+2)}\}$  is a Hamiltonian path between the vertices  $a_{1(1)}$  and  $a_{1(2h+2)}$  where  $s = t - (p + 1)$ ,  $p \geq 1$ .

## 4.2 Theorem

Let  $IG_N$  be the interleaver graph constructed from  $C(2n, 1, n)$  for  $n \geq 3$ . Then  $IG_N$ ,  $N \geq 6$  is 1-edge fault tolerant Hamiltonian  $t_e^*$  laceable for  $2 \leq t \leq \text{diam}(G)$ .

**Proof:** Let  $IG_N$  be an interleaver graph having cardinality  $N$ . Let  $V = V_1 \cup V_2$  be the set of vertices such that  $V_1 = \{a_{i(j)} : i = 1, 1 \leq j \leq N\}$ ,  $V_2 = \{a_{i(j)} : i = 2, 1 \leq j \leq N\}$ . Let  $E$  be the set of edges of  $IG_N$  such that  $E = \{b_k : 1 \leq k \leq 3N\} \cup E'$  where  $E' = \{a_{1(2)} - a_{2(2)}\}$  is the fault edge of  $IG_N$ .

**Case (i)**

In  $IG_N$ ,  $N = 4h + 2$  for  $h \geq 1$ , we have  $d(a_{1(1)}, a_{1(3)}) = 2$  and the path  $P : \{P_1 : a_{1(1)} - \phi[2] - a_{1(2)}\} \cup \{P_2 : a_{1(3)} - \phi[2h+2] - \{\Gamma[2] - H[x] - \phi[2] - H^{-1}[x]\}_{h-1} - \Gamma[2h+2] - S^{-1}[x] - \Gamma[2] - a_{2(2)}\} \cup \{P_3 : a_{1(2)} - L[x] - a_{2(2)}\}$  is a Hamiltonian path between the vertices  $a_{1(1)}$  and  $a_{1(3)}$ .

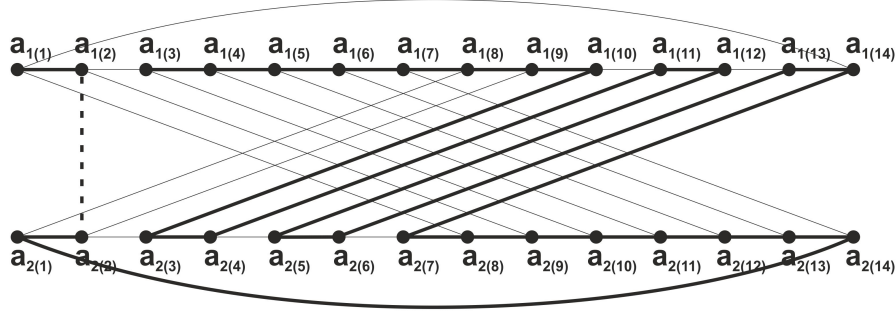


Figure 4 : Hamiltonian path between  $a_{1(1)}$  and  $a_{1(3)}$  in  $IG_{14}$ .

**Case (ii)**

In  $IG_N$ ,  $N = 4h + 2$  for  $h \geq 1$  we have  $d(a_{1(1)}, a_{1(2h+3)}) = t$  where  $4 \leq t \leq 2h + 2$  for even  $t$  and the path  $P : \{P_1 : a_{1(1)} - \phi[2] - a_{1(2)}\} \cup \{P_2 : a_{1(2h+3)} - \phi[2] - H^{-1}[x] - \{\Gamma[2] - H[x] - \phi[2] - H^{-1}[x]\}_{h-1} - \Gamma[4] - \{T^{-1}[x] - \phi[2] - T[x] - \Gamma[2]\}_{h-1} - T^{-1}[x] - \phi[2] - H^{-1}[x] - \Gamma[2] - a_{2(2)}\} \cup \{P_3 : a_{1(2)} - L[x] - a_{2(2)}\}$  is a Hamiltonian path between the vertices  $a_{1(1)}$  and  $a_{1(2h+3)}$ .

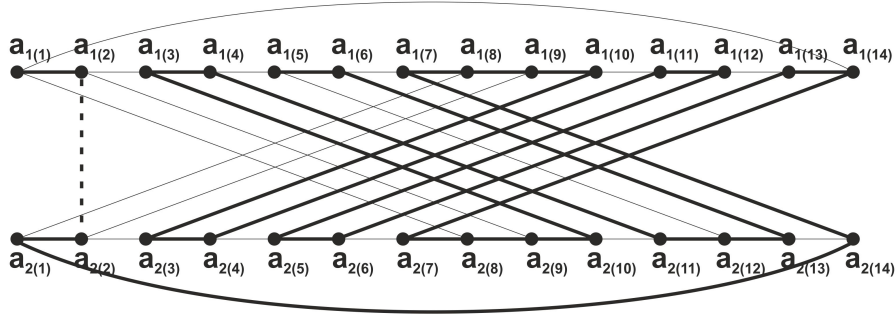


Figure 5 : Hamiltonian path between  $a_{1(1)}$  and  $a_{1(5)}$  in  $IG_{14}$ .

**Case (iii)**

In  $IG_N$ ,  $N = 4h + 6$  for  $h \geq 1$ , we have  $d(a_{1(1)}, a_{1(2h+3)}) = t$  such that  $4 \leq t \leq 2h + 2$  for even  $t$  and the path  $P : \{P_1 : a_{1(1)} - \phi[2] - a_{1(2)}\} \cup \{P_2 : a_{1(2h+3)} - \phi[4] - \{H^{-1}[x] - \Gamma[2] - H[x] - \phi[2]\}_h - H^{-1}[x] - \Gamma[4] - \{T^{-1}[x] - \phi[2] - T[x] - \Gamma[2]\}_h - S^{-1}[x] - \Gamma[2] - a_{2(2)}\} \cup \{P_3 : a_{1(2)} - L[x] - a_{2(2)}\}$  is a Hamiltonian path between the vertices  $a_{1(1)}$  and  $a_{1(2h+3)}$ .

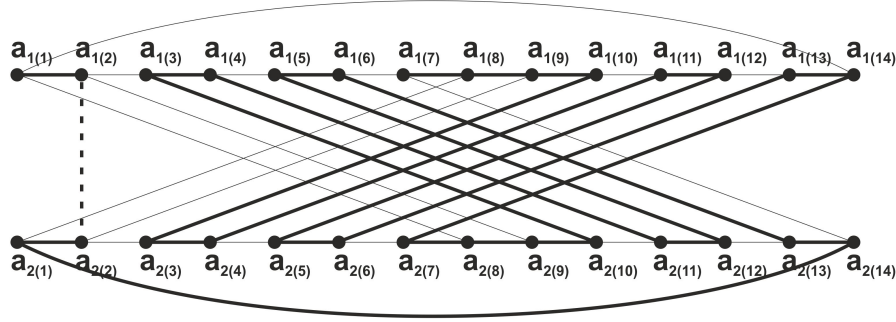


Figure 6 : Hamiltonian path between  $a_{1(1)}$  and  $a_{1(7)}$  in  $IG_{14}$ .

#### Case (iv)

In  $IG_N$ ,  $N = 4h + 10$  for  $h \geq 1$ , we have  $d(a_{1(1)}, a_{1(2h+3)}) = t$  such that  $4 \leq t \leq 2h + 2$  for even  $t$  and the path  $P : \{P_1 : a_{1(1)} - \phi[2] - a_{1(2)}\} \cup \{P_2 : a_{1(2h+3)} - \phi[2h] - T[x] - \Gamma^{-1}[2h + 1] - \{T^{-1}[x] - \phi^{-1}[2] - T[x] - \Gamma^{-1}[2]\}_{h-1} - T^{-1}[x] - \phi^{-1}[2] - T[x] - \Gamma^{-1}[4] - \{H[x] - \phi^{-1}[2] - H^{-1}[x] - \Gamma^{-1}[2]\}_{h+s} - H[x] - \phi^{-1}[4] - T[x] - S^{-1}[x] - \Gamma[2] - a_{2(2)}\} \cup \{P_3 : a_{1(2)} - L[x] - a_{2(2)}\}$  is a Hamiltonian path between the vertices  $a_{1(1)}$  and  $a_{1(2h+3)}$  where  $s = t - (p + 1)$ ,  $p \geq 1$ .

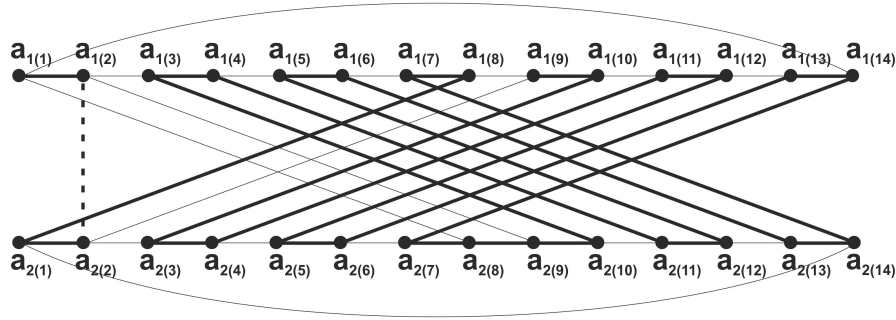


Figure 7 : Hamiltonian path between  $a_{1(1)}$  and  $a_{1(9)}$  in  $IG_{14}$ .

## Conclusion

The study of Hamiltonicity and Hamiltonian laceability has lot of significance in computer networks. In this paper, we have studied the Hamiltonian laceability properties of the interleaver graph  $IG_N$  for  $N \geq 6$ , constructed from the brick product graph  $C(2n, 1, n)$ . In particular, we have shown that the interleaver graph  $IG_N$ ,  $N \geq 6$  is Hamiltonian  $t_0^*$  laceable for  $1 \leq t \leq \text{diam}(G)$  and 1-edge fault tolerant Hamiltonian  $t_e^*$  laceable for  $2 \leq t \leq \text{diam}(G)$ . This concludes that the existence of Hamiltonian paths in such network suffice to solve many data communication problems.

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