# A Study of South East Corner Method and Use of Object Oriented Programming Model 

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#### Abstract

In this paper, the south east corner [ SEM ] procedure is successfully coded and tested via many randomly generated problem instances. Based on the results we can conclude that the correctness of the newly coded SEM is promising as compared with the previously coded one.


Keywords: Transportation problem, LPP, optimal solution, south east corner rule, object oriented programming.

## 1. INTRODUCTION

The term 'OR' was coined in 1940 by M. C. Closky \& T.ref then in a small town of Bawdsey in England. It is a science that came into existence in a military content. During world war II, the military management of UK called an Scientists from various disciplines \& organized them into teams to assist it in solving strategic \& tactical problems relating to air $\&$ land defence of the country.

The transportation problem is a special class of LPP that deals with shipping a product from multiple origins to multiple destinations. The objective of the transportation problem is to find a feasible way of transporting the shipments to meet demand of each destination that minimizes the total transportation cost while satisfying the supply \& demand constraints. The two basic steps of the transportation method are

Step 1: Determine the initial basic feasible solution
Step 2: Obtain the optimal solution using the solution obtained from step 1.
In this paper the corrected coding of SEM in $\mathrm{C}++$ is implemented. Then its correctness is verified via many randomly generated instances. The remainder of this paper is organized as follows :

Section II deals with the mathematical formulation of the transportation problem. In section III SEM is summarized. In section IV potential significance of the new object oriented
program of VAM is illustrated with a numerical example.
Finally, conclusion by highlighting the limitations and future research scope on the topic is made in section V.

## 2. MATHEMATICALFORMULATION OF THE TRANSPORTATION PROBLEM

A. In developing the LP model of the transportation problem the following notations are used
$a_{i}$ - Amounts to be shipped from shipping origin $i(a i \geq 0$.
$b_{j}$ - Amounts to be received at destination $\mathrm{j}(\mathrm{bj} \geq 0)$.
$\mathrm{c}_{\mathrm{ij}}$ - Shipping cost per unit from origin i to destination $\mathrm{j}(\mathrm{cij} \geq 0)$.
$\mathrm{x}_{\mathrm{ij}}$ - Amounts to be shipped from origin i to destination j to minimize the total cost
$\left(\mathrm{x}_{\mathrm{ij}} \geq 0\right)$

We assumed that the total amount shipped is equal to the total amount received, that is,

$$
\sum_{i=1}^{m} a_{i} \geq \sum_{i=1}^{m} b_{j}
$$

B. Transportation problem

$$
\begin{aligned}
& \operatorname{Min} \sum_{i=1}^{m} \sum_{j=1}^{n} c_{i j} x_{i j} \\
& \text { Subject to } \sum_{j=1}^{n} x_{i j} \leq \mathrm{a}_{\mathrm{i}}, \mathrm{i}=1,2, \ldots, \mathrm{~m} \\
& \sum_{i=1}^{m} x_{i j} \leq \mathrm{b}_{\mathrm{j}}, \mathrm{j}=1,2, \ldots, \mathrm{n}, \text { where } x_{i j} \geq 0 \forall \mathrm{i}, \mathrm{j} .
\end{aligned}
$$

Feasible solution : A set of non negative values $x_{i j}, \mathrm{i}=1,2, \ldots, \mathrm{n}$ and $\mathrm{j}=1,2, \ldots, \mathrm{~m}$ that satisfies the constraints is called a feasible solution to the transportation problem .

Optimal solution : A feasible solution is said to be optimal if it minimizes the total transportation cost .
Non degenerate basic feasible solution: A basic feasible solution to a ( $\mathrm{m} \times \mathrm{n}$ ) transportation problem that contains exactly $m+n-1$ allocations in independent positions.

Degenerate basic feasible solution: A basic feasible solution that contains less that $m+n-1$ non negative allocations.

Balanced and Unbalanced Transportation problem: A Transportation problem is said to be balanced if the total supply from all sources equals the total demand in the destinations and is called unbalanced otherwise.

Thus, for a balanced problem, $\sum_{i=1}^{m} a_{i}=\sum_{i=1}^{m} b_{j}$ and for unbalanced problem, $\sum_{i=1}^{m} a_{i} \neq \sum_{i=1}^{m} b_{j}$

## 3. SOUTH EAST CORNER RULE

The steps involved in VAM in producing the initial feasible solution are described below :
This method starts at the south east corner cell (route) of the table variable $\left(\mathrm{x}_{43}\right)$.
Step 1: Construct the transportation table for the given TPP.
Step 2: Allocate as much as possible to the selected cell and adjust the associated amounts of supply and demand by subtracting the allocated amount.

Step 3: Cross out the row or column with zero supply or demand to indicate that no further assignments can be made in that row or column. I both a row and column net to zero simultaneously, cross out one only and leave a zero supply
(demand in the uncrossed out row - column).
Step 4: If exactly one row or column is left uncrossed out, stop. Otherwise, move to the cell to the right if a column has just been crossed out. Go to step 2 [8].

## 4. Object oriented program code for South east corner method for initial basic feasible solution: **************************************************

## SOUTH EAST CORNER

```
**************************************************/
#include<stdio.h>
#include<conio.h>
void main()
{
    int sn,dn,i,j,ss=0,ds=0,sum=0;
    int sup[10],dem[10];
    int a[10][10],c[10][10];
    clrscr();
    //getting no.of supply & demand.
    printf("Enter the num of supply:");
    scanf("%d",&sn);
    printf("Enter the num of demand:");
    scanf("%d",&dn);
    //clearing values in array.
    for(i=0,j=0;i<10,j<10;i++,j++)
    {
            sup[i]=0;
            dem[j]=0;
    }
```

    //input supply values \& calculate sum of supply.
    for \((\mathrm{i}=0 ; \mathrm{i}<\mathrm{sn} ; \mathrm{i}++\) )
    \{
printf("\n Enter supply value sup[\%d]:",i);
scanf("\%d",\&sup[i]);
ss=ss+sup[i];
\}
//input demand values \& calculate sum of demand.

$$
\operatorname{for}(\mathrm{j}=0 ; \mathrm{j}<\mathrm{dn} ; \mathrm{j}++)
$$

\{
printf("\n Enter demand value dem[\%d]:",j);
scanf("\%d",\&dem[j]);
ds=ds+dem[j];
\}
if(ss!=ds)
\{
printf("\n unbalanced problem..");
getch();
exit(0);
\}
//input transportation cost.
printf("\n Enter array values:");

```
for(i=0;i<sn;i++)
```

\{

```
    for(j=0;j<dn;j++)
    {
        scanf("%d",&a[i][j]);
```

```
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    }
}
//clearing cost array.
for(i=0;i<10;i++)
{
    for(j=0;j<10;j++)
    {
        c[i][j]=0;
    }
}
//calculation.
for(i=0,j=dn-1;i<sn,j>=0;)
{
    if(i==sn||j==-1)
            goto L1;
    if(dem[j]==sup[i]) //Checking demand=supply
    {
        c[i][j]=dem[j]*a[i][j];
        printf("\n c[%d][%d]=%d and sup=%d",i,j,c[i][j],sup[i]);
        j--;
        i++;
        }
        else if(dem[j]<sup[i])//checking demand< supply
        {
        c[i][j]=dem[j]*a[i][j];
```

            sup[i]=sup[i]-dem[j];
            printf("\n c[%d][%d]=%d and sup=%d",i,j,c[i][j],sup[i]);
            dem[j]=0;
            j--;
                }
            else //demand >= supply
            {
            c[i][j]=sup[i]*a[i][j];
            dem[j]=dem[j]-sup[i];
            printf("\n c[%d][%d]=%d and dem=%d",i,j,c[i][j],dem[j]);
            i++;
            }
    }
    L1:
for(i=0;i<sn;i++)
for(j=0;j<dn;j++)
sum=sum+c[i][j];
printf("\n\n sum of transportation cost = %d",sum);
getch();
}

```

\section*{Output :}

Enter the num of supply:4
Enter the num of demand:5
Enter supply value \(\sup [0]: 800\)
Enter supply value sup[1]:500

Enter supply value sup[2]:900
Enter supply value sup[3]:300
Enter demand value dem[0]:400
Enter demand value dem[1]:400
Enter demand value dem[2]:500
Enter demand value dem[3]:400
Enter demand value dem[4]:800
Enter array values:5 8663
\[
47765
\]

84664
00000
\(c[0][4]=2400\) and sup \(=800\)
\(c[1][3]=2400\) and sup \(=100\)
\(c[1][2]=700\) and dem=400
c[2][2]=2400 and sup=500
\(c[2][1]=1600\) and \(\sup =100\)
\(c[2][0]=800\) and dem=300
\(c[3][0]=0\) and \(\sup =300\)
sum of transportation cost \(=10300\)

\section*{5. Conclusion :}

The optimal solution obtained in this present investigation shows much more closeness with initial basic feasible solution obtained by South east corner rule. The comparison of optimal solution have been made with other methods of finding initial solutions and observe that South east corner method give the better initial feasible solutions which are closer to optimal solution. The objected oriented programs using c++ have been developed . This shows that the computed results tally with the results obtained \(\mathrm{c}++\) programming. . Object oriented program code for said programs is given for better understanding.

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